



Information Technology Engineering

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Voice and Video over IP

MULTIMEDIA NETWORKING

Slides derived from those available on the Web site of the book
“Computer Networking”, by Kurose and Ross, PEARSON

Multimedia networking: outline

7.1 multimedia networking applications

7.2 streaming *stored* video

7.3 **voice-over-IP**

7.4 protocols for *real-time* conversational applications

7.5 network support for multimedia

Voice-over-IP (VoIP)

- ❖ *VoIP end-end-delay requirement*: needed to maintain “conversational” aspect
 - higher delays noticeable, impair interactivity
 - < 150 msec: good
 - > 400 msec bad
 - includes application-level (packetization, playout), network delays
- ❖ *session initialization*: how does callee advertise IP address, port number, encoding algorithms?
- ❖ *value-added services*: call forwarding, screening, recording
- ❖ *emergency services*: 911

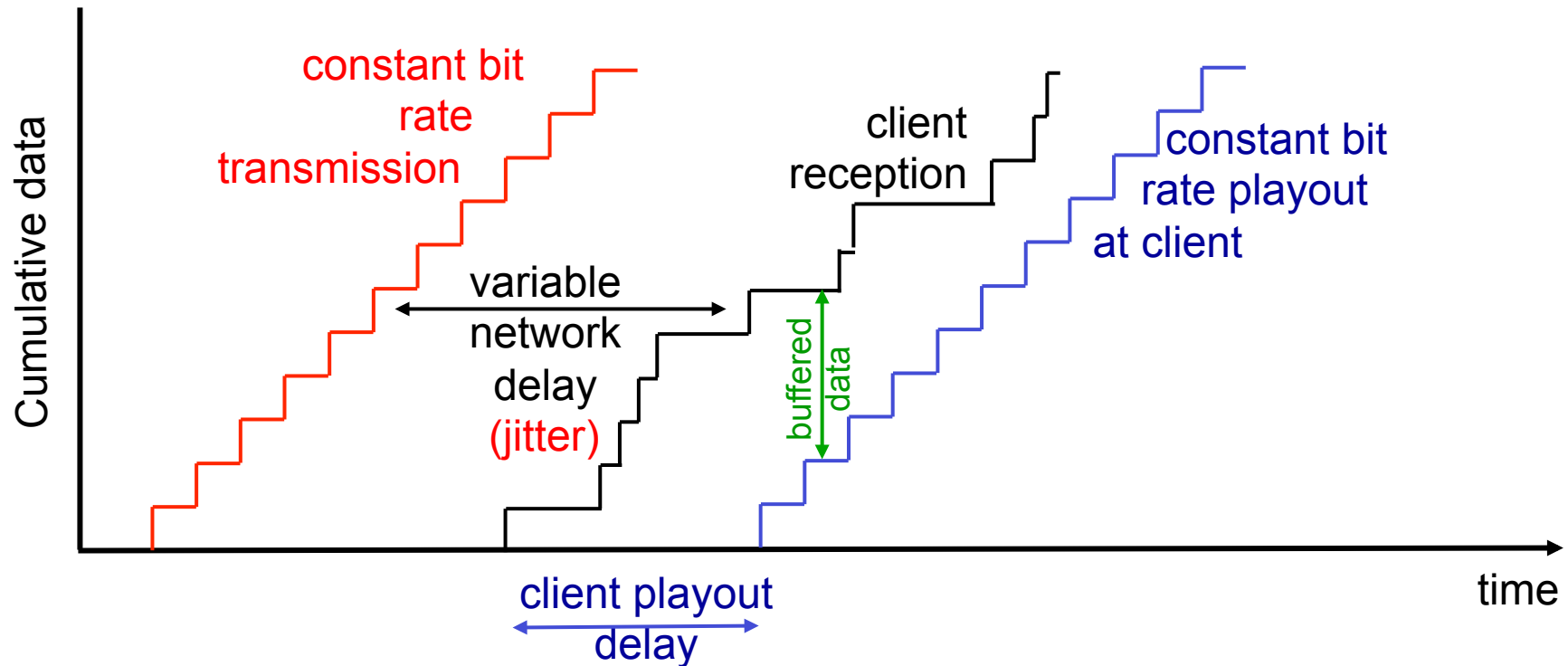
VoIP characteristics

- ❖ Speaker's audio: alternating talk spurts, silent periods.
 - 64 kbps during talk spurt
 - pkts generated only during talk spurts
 - 20 msec chunks at 8 Kbytes/sec: 160 bytes of data
- ❖ application-layer header added to each chunk
- ❖ chunk+header encapsulated into UDP or TCP segment
- ❖ application sends segment into socket every 20 msec during talkspurt

VoIP: packet loss, delay

- ❖ *network loss*: IP datagram lost due to network congestion (router buffer overflow)
- ❖ *delay loss*: IP datagram arrives too late for playout at receiver
 - delays: processing, queueing in network; end-system (sender, receiver) delays
 - typical maximum tolerable delay: 400 ms
- ❖ *loss tolerance*: depending on voice encoding, loss concealment, packet loss rates between 1% and 10% can be tolerated

Delay jitter



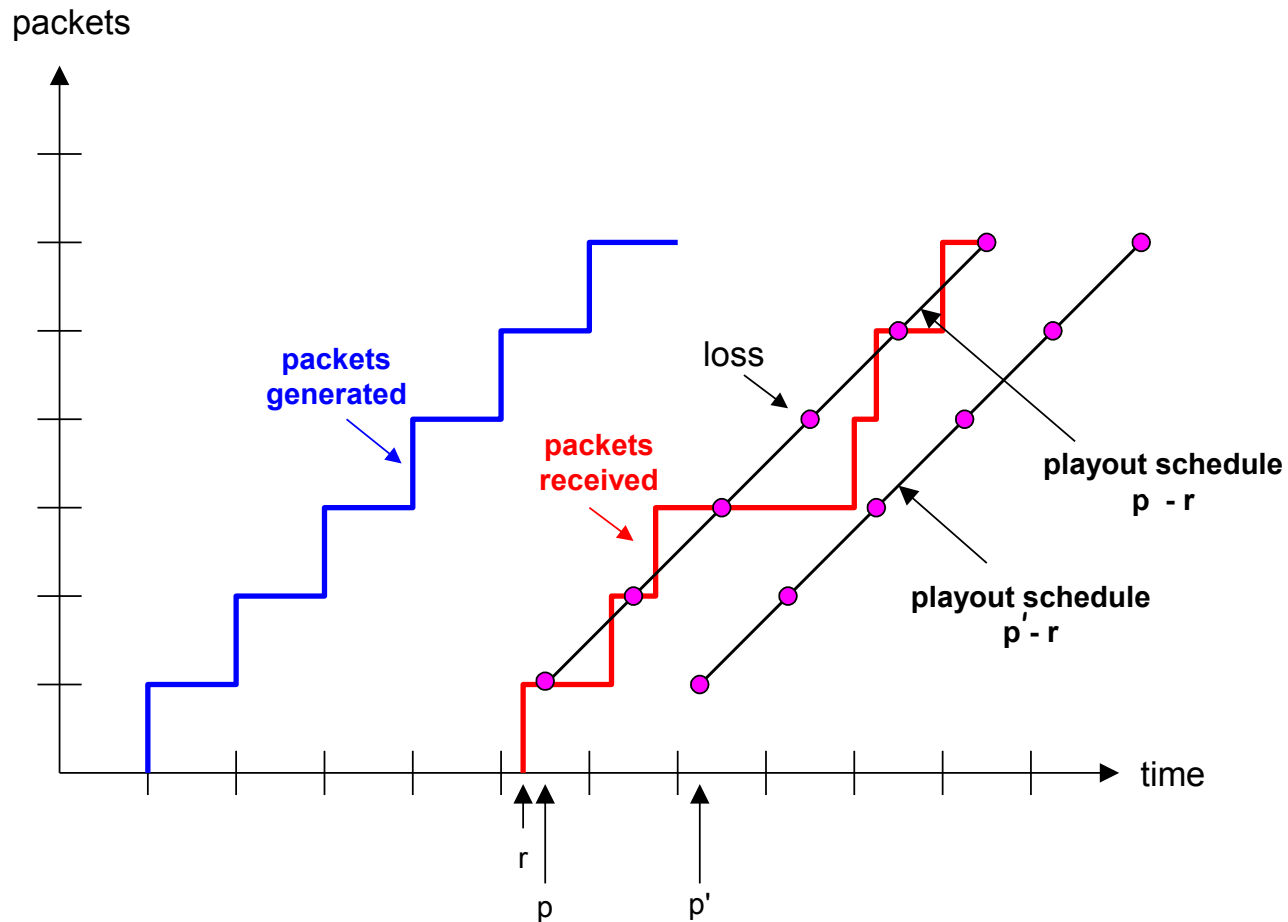
- ❖ end-to-end delays of two consecutive packets: difference can be more or less than 20 msec (transmission time difference)

VoIP: fixed playout delay

- ❖ receiver attempts to playout each chunk exactly q msec after chunk was generated.
 - chunk has time stamp t : play out chunk at $t+q$
 - chunk arrives after $t+q$: data arrives too late for playout: data “lost”
- ❖ tradeoff in choosing q :
 - *large q : less packet loss*
 - *small q : better interactive experience*

Fixed Playout Delay

- Sender generates packets every 20 msec during talk spurt.
- First packet received at time r
- First playout schedule: begins at p
- Second playout schedule: begins at p'



Adaptive playout delay (I)

- ❖ *goal*: low playout delay, low late loss rate
- ❖ *approach*: adaptive playout delay adjustment:
 - estimate network delay, adjust playout delay at beginning of each talk spurt
 - silent periods compressed and elongated
 - chunks still played out every 20 msec during talk spurt
- ❖ adaptively estimate packet delay: (EWMA - exponentially weighted moving average, **recall TCP RTT estimate**):

$$d_i = (1-\alpha)d_{i-1} + \alpha (r_i - t_i)$$

delay estimate after ith packet *small constant, e.g. 0.1* *time received - time sent (timestamp)*
measured delay of ith packet

Adaptive playout delay (2)

- ❖ also useful to estimate average deviation of delay, v_i :

$$v_i = (1-\beta)v_{i-1} + \beta |r_i - t_i - d_i|$$

- ❖ estimates d_i , v_i calculated for every received packet, but used only at start of talk spurt

- ❖ for first packet in talk spurt, playout time is:

$$\text{playout-time}_i = t_i + d_i + Kv_i$$

remaining packets in talkspurt are played out periodically

Adaptive playout delay (3)

Q: How does receiver determine whether packet is first in a talkspurt?

- ❖ if no loss, receiver looks at successive timestamps
 - difference of successive stamps > 20 msec --> talk spurt begins.
- ❖ with loss possible, receiver must look at both time stamps and sequence numbers
 - difference of successive stamps > 20 msec *and* sequence numbers without gaps --> talk spurt begins.

VoiP: recovery from packet loss (I)

Challenge: recover from packet loss given small tolerable delay between original transmission and playout

- ❖ each ACK/NAK takes \sim one RTT
- ❖ alternative: *Forward Error Correction (FEC)*
 - send enough bits to allow recovery without retransmission (recall two-dimensional parity in Ch. 5)

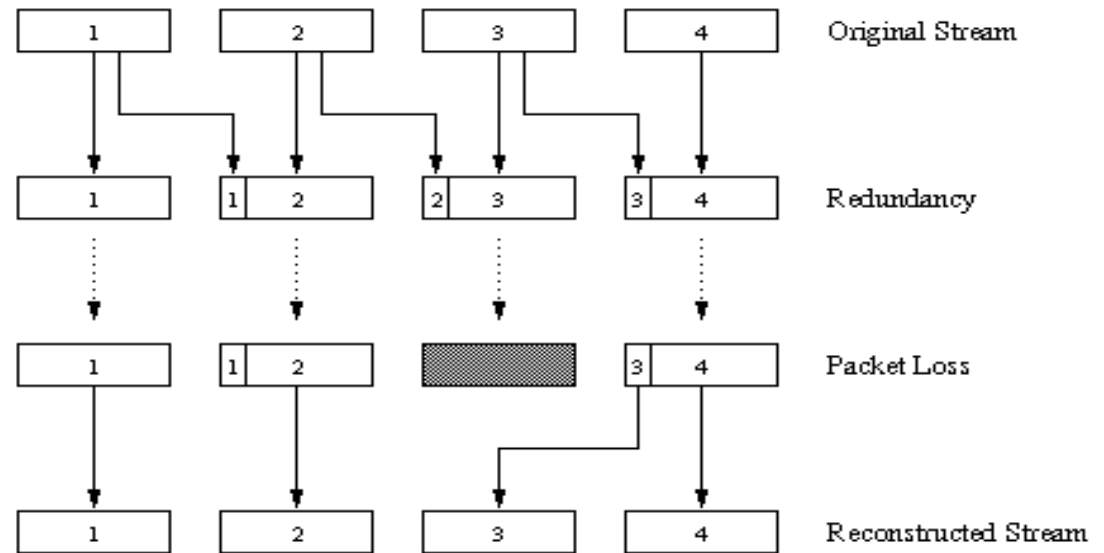
Simple FEC

- ❖ for every group of n chunks, create redundant chunk by exclusive OR-ing n original chunks
- ❖ send $n+1$ chunks, increasing bandwidth by factor $1/n$
- ❖ can reconstruct original n chunks if at most one lost chunk from $n+1$ chunks, with playout delay

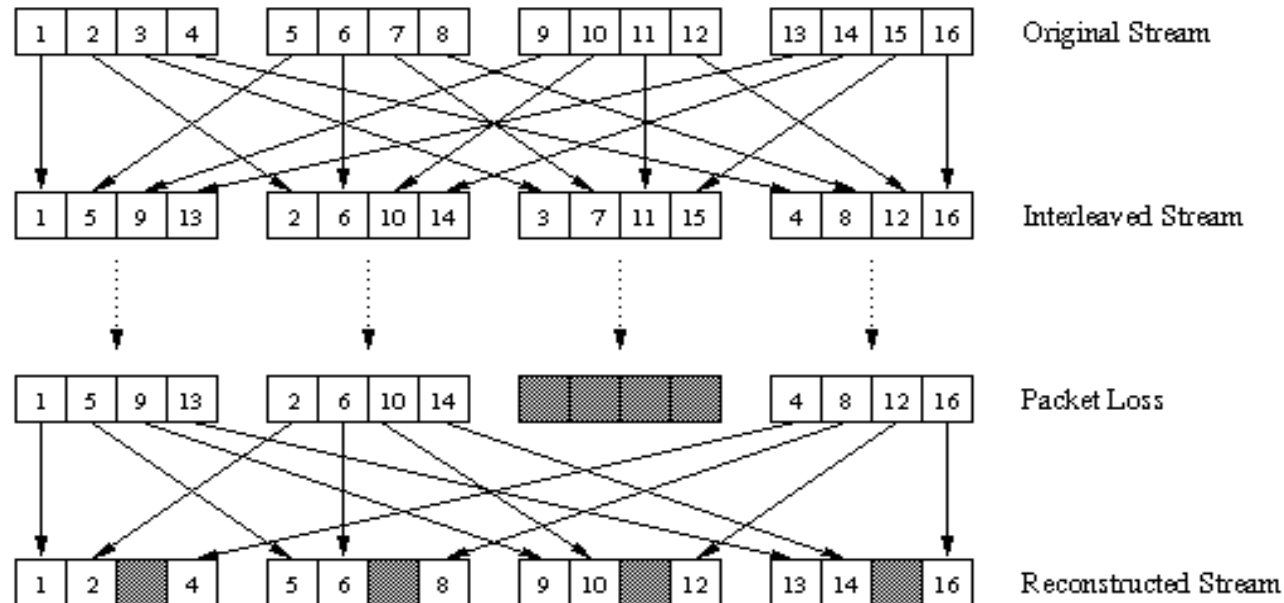
VoiP: recovery from packet loss (2)

another FEC scheme:

- ❖ “piggyback lower quality stream”
- ❖ send lower resolution audio stream as redundant information
- ❖ e.g., nominal stream PCM at 64 kbps and redundant stream GSM at 13 kbps
- ❖ non-consecutive loss: receiver can conceal loss
- ❖ generalization: can also append (n-1)st and (n-2)nd low-bit rate chunk



VoiP: recovery from packet loss (3)

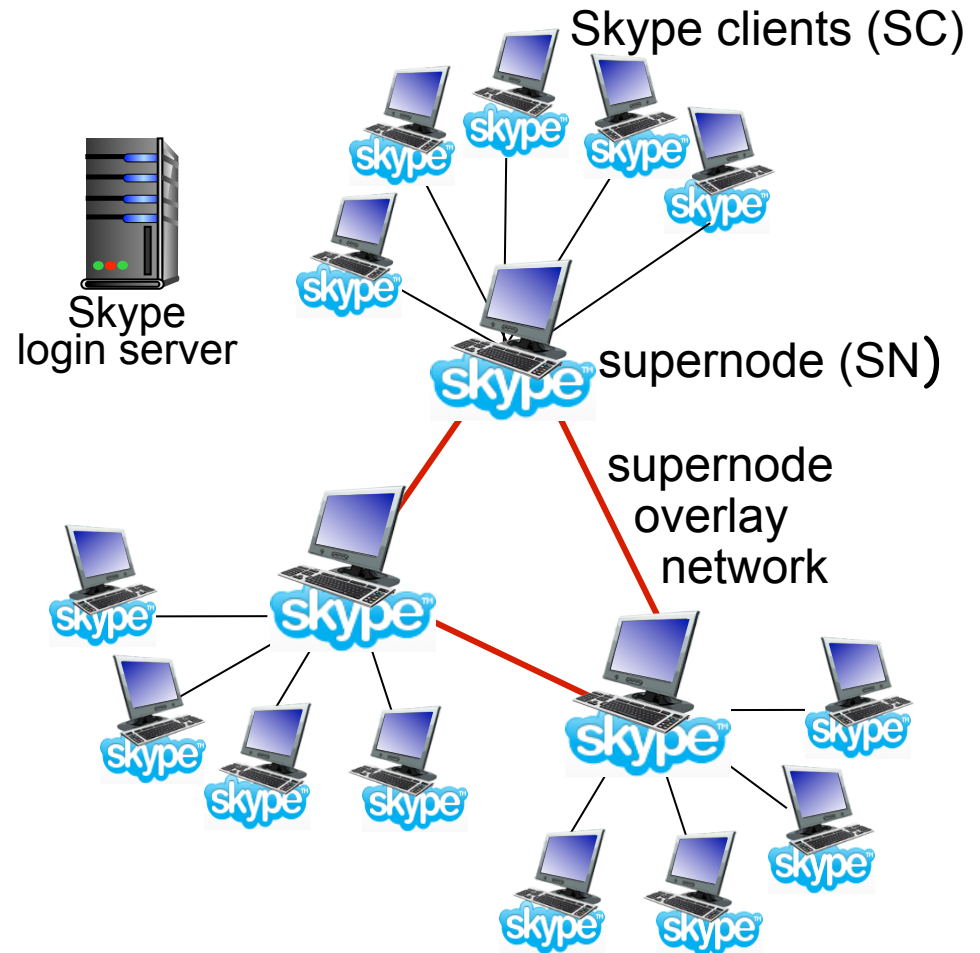


interleaving to conceal loss:

- ❖ audio chunks divided into smaller units, e.g. four 5 msec units per 20 msec audio chunk
- ❖ packet contains small units from different chunks
- ❖ if packet lost, still have *most* of every original chunk
- ❖ no redundancy overhead, but increases playout delay

Voice-over-IP: Skype

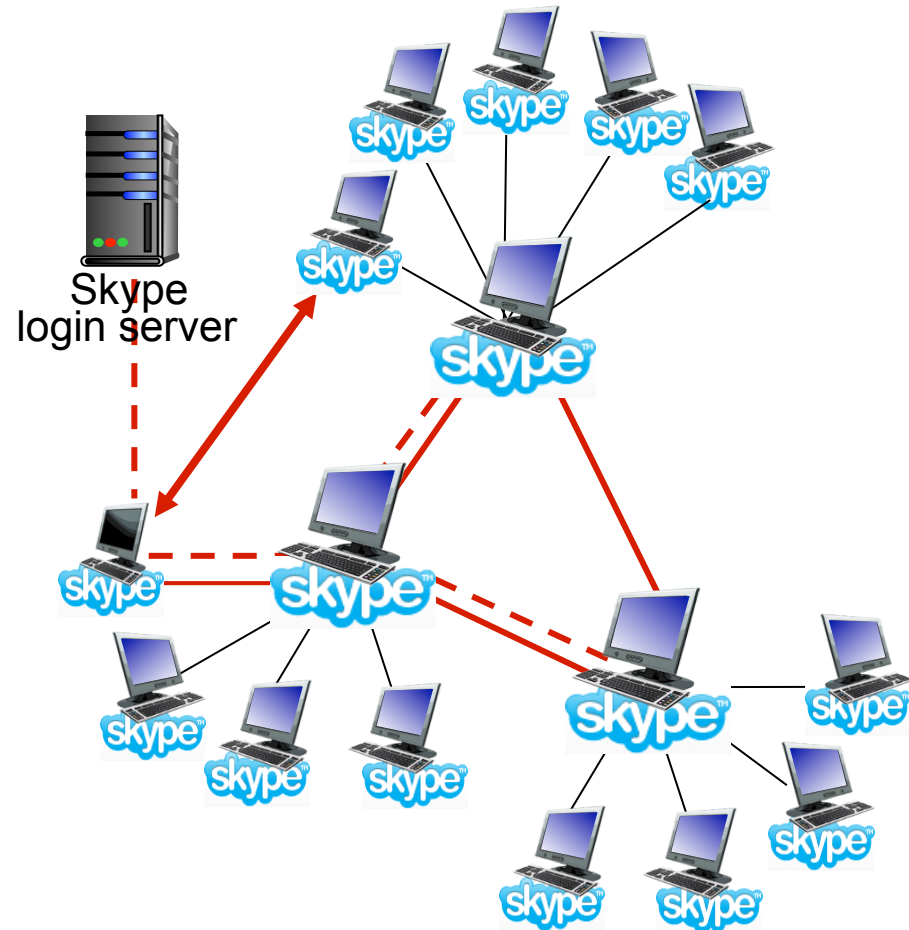
- ❖ proprietary application-layer protocol (inferred via reverse engineering)
 - encrypted msgs
- ❖ P2P components:
 - **clients**: skype peers connect directly to each other for VoIP call
 - **super nodes (SN)**: skype peers with special functions
 - **overlay network**: among SNs to locate SCs
 - **login server**



P2P voice-over-IP: skype

skype client operation:

1. joins skype network by contacting SN (IP address cached) using TCP
2. logs-in (username, password) to centralized skype login server
3. obtains IP address for callee from SN, SN overlay
 - or client buddy list
4. initiate call directly to callee



Skype: Peers as Relays

- ❖ **problem:** both Alice, Bob are behind “NATs”
 - NAT prevents outside peer from initiating connection to insider peer
 - inside peer *can* initiate connection to outside
- ❖ **relay solution:** Alice, Bob maintain open connection to their SNs
 - Alice signals her SN to connect to Bob
 - Alice’s SN connects to Bob’s SN
 - Bob’s SN connects to Bob over open connection Bob initially initiated to his SN

